

Decycling Number of Circular Graphs*

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Abstract The lower and upper bounds on the decycling number of circular graph $C(n, k)$ where $k \leq \lfloor \frac{n}{2} \rfloor$ of order n are obtained. The explicit expressions of that of some classes of graphs are presented.

Keywords circular graph; decycling number; independent set

1 Introduction: the decycling number of graphs

It is well known that the cycle rank of a graph is the minimum number of edges whose removal eliminates all cycles in the graph. The parameter has a simple expression. That is, if G is a graph with p vertices, q edges and k components, then the cycle rank $\beta(G) = q - p + k$. It is an important invariant to characterize a graph. The corresponding problem of removing vertices does not have such a simple solution. It is quite difficult even for some elementary graph.

Let $G(V, E)$ be a graph. If $S \subseteq V(G)$ and $G - S$ is acyclic, then S is said to be a *decycling set* of G . The minimum order of decycling set is called the *decycling number* of G and is denoted by $\nabla(G)$. A decycling set of this order is called a ∇ -set. It was shown[4] that determining the decycling number of an arbitrary graph is NP-complete. The results on the decycling number of several classes of simply defined graphs can be seen in [1-2,7].

For the basic terminologies and notations, we refer the reader to [3].

2 The decycling number of circular graph

Let n and l be two positive integers with $n \geq 2l$. For any two numbers i, j where $1 \leq i, j \leq n$, define function

$$\chi(i-j) = \begin{cases} |i-j|, & \text{if } |i-j| \leq \frac{n}{2} \\ n - |i-j|, & \text{otherwise.} \end{cases}$$

A circular graph $G = C(n, l)$ of order n is one spanned by a n -circuit $C_n = (1, 2, \dots, n)$ together with the chords $(i, j) \in E(G)$ iff $\chi(j-i) = l$ ($l > 1$). Circular graphs are very

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useful for their own sake. In literature [5,6], the crossing number and some topological results of circular graphs were shown. In this paper the decycling number of circular graph is discussed.

Theorem 1 $\lceil \frac{n+1}{3} \rceil \leq \nabla(C(n,l)) \leq \lceil \frac{n}{2} \rceil$ where $n > 2l$ and $l > 1$.

Proof. It is easy to see that graph $C(n,l)$ is 4-regular. Firstly we prove that $\nabla(C(n,l)) \geq \lceil \frac{n+1}{3} \rceil$. Suppose S is a decycling set of graph $C(n,l)$ with order m . To make graph $C(n,l) - S$ acyclic, at most $n - m - 1$ edges are allowed. That is, at least $2n - (n - m - 1) = n + m + 1$ edges should be removed. Let $E(S) = \{e = (u,v) \mid u \text{ or } v \in S\}$. Then $|E(S)| \leq 4m$ and $|E(S)| = 4m$ if and only if S is an independent set. Then $4m \geq n + m + 1$ is obtained. It follows that $m \geq \lceil \frac{n+1}{3} \rceil$.

Then we prove that $\nabla(C(n,l)) \leq \lceil \frac{n}{2} \rceil$.

Case 1 n is even and l is odd. Let $S = \{x \mid x \text{ is even, } x \leq n\}$. $\forall y \in V(C(n,l)) - S$, if $(y,z) \in E(C(n,l))$, z is even. So $z \in S$. That is to say that $C(n,l) - S$ is a set of $\frac{n}{2}$ isolated vertices. So $\nabla(C(n,l)) \leq \frac{n}{2}$.

Case 2 n is even and l is even. If $n = 4k$ let

$$S = \{x \mid x \text{ is even and } x \leq \frac{n}{2}\} \cup \{y \mid y \text{ is odd and } \frac{n}{2} < x < n\}.$$

If $n = 4k + 2$ let

$$S = \{x \mid x \text{ is even and } x < \frac{n}{2}\} \cup \{y \mid y \text{ is odd and } \frac{n}{2} < x < n\} \cup \{\frac{n}{2} + 1\}.$$

All of the edges $(i, i+1)$, $i = 1, 2, \dots, n-1$, are not in $G - S$. Now we prove that S is a decycling set of G . Assume C is a cycle in $G - S$ and C does not contain the edge $(1, n)$, then it is composed of the edges $(i, i+l)$. Without loss of generality, we can suppose that $C = \{i, i+l, i+2l, \dots, i+xl\}$ where the numbers are taken modulo n and $i < \frac{n}{2}$. We can see that $i + xl + l - n = i$. Then i is odd, $i + xl$ is odd too but $i + xl$ belongs to the set S . So such a circuit does not exist.

Assume C is a cycle in $G - S$ and C contains the edge $(1, n)$, since $1 + n - l \in S$, $n + l \equiv l \in S$, the circuit C is $\{1 + xl, \dots, 1 + 2l, 1 + l, 1, n, n - l, n - 2l, \dots, n - yl\}$ and $1 + xl + l \equiv n - yl$ which is impossible since $1 + (x+1)l$ is odd and $n - yl$ is even. Then graph $C(n,l) - S$ is acyclic. So $\nabla(C(n,l)) \leq \lceil \frac{n}{2} \rceil$.

Case 3 n is odd and l is odd. Let $S = \{x \mid x \text{ is even, } x \leq n\} \cup \{n\}$. In graph $C(n,l) - S$, given an odd number y ($l \leq y < n - l$), it is an isolated vertex. For an odd number y ($0 < y < l$), it is adjacent to only one vertex $n + y - l$. That is, it is an articulate vertex. For the same reason, the vertex y ($n - l < y < n$) is adjacent to only one vertex $y + l - n$. Graph $C(n,l) - S$ is acyclic. So $\nabla(C(n,l)) \leq \lfloor \frac{n}{2} \rfloor + 1 = \lceil \frac{n}{2} \rceil$.

Case 4 n is odd and l is even. Let $S = \{x \mid x \text{ is even, } x \leq n\} \cup \{n\}$. For any $y \in V(C(n,l)) - S$, at least two edges of E_y are removed. If there is a circuit C in graph $(C(n,l)) - S$, C must be composed of such edges as $(i, i+l)$ where the numbers are modulo n . Such an edge (x,y) ($n - l \leq x \leq n$, $0 \leq y \leq l$) is inevitable which contradicts the fact that x and y should be odd at the same time. So $C(n,l) - S$ is acyclic. Then $\nabla(C(n,l)) \leq \lceil \frac{n}{2} \rceil$.

□

From the proof of this theorem the following corollary is direct:

Corollary 2 Suppose $n \equiv 2 \pmod{3}$ and $\nabla(C(n, l)) = |S| = \frac{n+1}{3}$ for certain l where S is a decycling set and $n \neq 2l$. Then the decycling set is an independent set and the graph $G - S$ is a tree (connected and acyclic).

Lemma 3[2] If G and H are homomorphic graphs, then $\nabla(G) = \nabla(H)$.

In the following paper, we determine the explicit expression of ∇ -set for certain circular graphs. At first if $n = 2l$ then graph $C(n, l)$ is 3-regular. We discuss the decycling set of $C(n, l)$.

Lemma 4 $\nabla(C(n, l)) \geq \lceil \frac{l+1}{2} \rceil$ where $n = 2l$.

Proof. Suppose S is a decycling set with $|S| = m$ of graph $C(n, l)$. The size of graph $C(n, l) - S$ is at most $l - m - 1$. That is to say, at least $l + m + 1$ edges should be eliminated. Then we get $3m \geq l + m + 1$ which follows that $m \geq \lceil \frac{l+1}{2} \rceil$. \square

Theorem 5 $\nabla(C(n, l)) = \lceil \frac{l+1}{2} \rceil$ where $l \geq 2$ and $n = 2l$.

Proof. Easily to see that graph $C(4, 2)$ is homomorphic to the complete graph K_4 . From [6], we know that $\nabla(C(4, 2)) = 2$.

Suppose A is a set of edges of graph G . Then $\nabla(G - A) \leq \nabla(G)$. For any edge $(i, j) \in E(C(n, l))$ where $|i - j| \neq 1$ graph $C(n, n) - (i, j)$ is homomorphic to graph $C(n - 2, l - 1)$. From Lemma 3, $\nabla(C(n - 2, l - 1)) \leq \nabla(C(n, l)) \leq \nabla(C(n + 2, l + 1))$.

By induction on the number l . When $l = 3$, the circular graph $C(6, 3)$ is isomorphic to $K_{3,3}$, so $\nabla(C(6, 3)) = 2 = \lceil \frac{3+1}{2} \rceil$. If $l = 4$, from Lemma 4, $\nabla(C(8, 4)) \geq 3$. Suppose $S = \{1, 3, 5\}$, $C(8, 4)$ is acyclic. Then $\nabla(C(8, 4)) = 3 = \lceil \frac{4+1}{2} \rceil$.

Suppose when $nl = k$ (k is even) this theorem holds. We can get that $\frac{k}{2} + 1 \leq \nabla(C(2k + 2, k + 1))$. Let $S = \{i \mid i \text{ is odd, and } i \leq k + 1\}$. The induced graph $C(2k + 2, k + 1) - S$ is a tree. So $\nabla(C(2k + 2, k + 1)) = \frac{k}{2} + 1 = \lceil \frac{k+2}{2} \rceil$.

Suppose when $n = k$ (k is odd) this theorem holds. It is got that $\frac{k-1}{2} + 1 \leq \nabla(C(2k + 2, k + 1))$. On the other hand, from Lemma 4 $\nabla(C(2k + 2, k + 1)) \geq \frac{k-1}{2} + 2$. Let $S = \{i \mid i \text{ is odd, and } i \leq k + 1\}$. The induced graph $C(2k + 2, k + 1) - S$ is a tree. So $\nabla(C(2k + 2, k + 1)) = \frac{k+1}{2} + 1 = \lceil \frac{k+2}{2} \rceil$.

This theorem follows. \square

Most circular graphs $C(n, l)$ is 4-regular. Then we consider 4-regular circular graph. If $n = 2$ we have the following result:

Theorem 6

$$\nabla(C(n, 2)) = \begin{cases} \lceil \frac{n+1}{3} \rceil + 1, & \text{if } n \equiv 2, 5 \pmod{6} \\ \lceil \frac{n+1}{3} \rceil, & \text{otherwise.} \end{cases}$$

where $n \geq 5$.

Proof. When $n = 5$, $C(5, 2)$ is isomorphic to K_5 . And $\nabla(C(5, 2)) = 3 = \lceil \frac{5+1}{3} \rceil + 1$.

Then we prove the case when $n \geq 6$. First from Theorem 1, we know $\nabla(C(n, 2)) \geq \lceil \frac{n+1}{3} \rceil$.

Case 1 $n = 6k$ where k is a positive integer. On one hand, $\nabla(C(6k, 2)) \geq 2k + 1$. Suppose $S = \{3i \mid i = 1, 2, \dots, 2k\} \cup \{1\}$. Easy to see that $C(6k, 2) - S$ is a path of length $4k - 2$. Then $\nabla(C(6k, 2)) = 2k + 1 = \lceil \frac{n+1}{3} \rceil$.

Case 2 $n = 6k + 2$ where k is a positive integer. It is easy to see that $\nabla(C(6k, 2)) \geq 2k + 1$. We say that $\nabla(C(6k + 2, 2)) \geq 2k + 2$. Otherwise suppose that S is a decycling

set with order $2k + 1$ of graph $C(6k + 2, 2)$. From Corollary 2, $|E(S)| = 4k + 4$. Select any $2k + 1$ vertices from $6k + 2$ vertices. Since $\frac{6k+2}{2k+1} < 3$, the occurrence of such pairs of vertices as i and $i + 1$ or i and $i + 2$ is inevitable. So $|E(S)| = 4k + 4$ is impossible. Then we get that $\nabla(C(6k, 2)) \geq 2k + 2$. On the other hand, let $S = \bigcup_{i=1}^{2k} \{3i\} \cup \{6k + 1, 6k + 2\}$. The graph $C(6k, 2) - S$ is a path of length $4k - 1$. Then $\nabla(C(6k + 2, 2)) = 2k + 2 = \lceil \frac{6k+3}{3} \rceil + 1$.

Case 3 $n = 6k + 4$ where k is a positive integer. On one hand, $\nabla(C(6k + 4, 2)) \geq 2k + 2$. On the other hand let $S = \bigcup_{i=1}^{2k+1} \{3i\} \cup \{6k + 4\}$. The graph $C(6k, 2) - S$ is a path of length $4k + 1$. So $\nabla(C(6k + 4, 2)) = 2k + 2 = \lceil \frac{6k+5}{3} \rceil$.

Case 4 $n = 6k + 1$ where k is a positive integer. From Theorem 1, $\nabla(C(6k + 1, 2)) \geq 2k + 1$. Then let $S = \bigcup_{i=1}^{2k} \{3i\} \cup \{6k + 1\}$. One can see that graph $C(6k + 1, 2) - S$ is a path of length $4k - 1$. So $\nabla(C(6k + 1, 2)) = 2k + 1 = \lceil \frac{6k+2}{3} \rceil$.

Case 5 $n = 6k + 3$ where k is a positive integer. On one hand, From Theorem 1 $\nabla(C(6k + 3, 2)) \geq 2k + 2$. On the other hand let $S = \bigcup_{i=1}^{2k+1} \{3i\} \cup \{1\}$. The graph $C(6k + 3, 2) - S$ is a path of length $4k$. So $\nabla(C(6k + 3, 2)) = 2k + 2 = \lceil \frac{6k+4}{3} \rceil$.

Case 6 $n = 6k + 5$ where k is a positive integer. On one hand, $\nabla(C(6k + 5, 2)) \geq 2k + 2$. In graph $C(6k + 5, 2)$, any $2k + 2$ vertices can be incident to at most $8k + 7$ edges since the occurrence of such pair of vertices as $(i, i + 1)$ or $(i, i + 2)$ is inescapable. So $\nabla(C(6k + 5, 2)) \geq 2k + 3$. Let $S = \bigcup_{i=1}^{2k+1} \{3i\} \cup \{6k + 4, 6k + 5\}$. The graph $C(6k + 5, 2) - S$ is acyclic. Then $\nabla(C(6k + 5, 2)) = 2k + 3 = \lceil \frac{6k+6}{3} \rceil + 1$.

This theorem is found. □

Theorem 7

$$\nabla(C(n, 3)) = \begin{cases} \lceil \frac{n+1}{3} \rceil + 1, & \text{if } n = 3k + 2 \text{ and } k \text{ is odd,} \\ \lceil \frac{n+1}{3} \rceil, & \text{otherwise.} \end{cases}$$

where $n \geq 7$.

Proof First from Theorem 1, it is known that $\nabla(C(n, 3)) \geq \lceil \frac{n+1}{3} \rceil$.

Case 1 $n = 3k$ where k is a positive integer. Suppose that $S = \{a_t, n - 1\}$ where $a_1 = 1$, $a_t < n$ and

$$a_{t+1} = \begin{cases} a_t + 2, & \text{if } t \text{ is odd,} \\ a_t + 4, & \text{if } t \text{ is even.} \end{cases}$$

There is no cycle in graph $C(3k, 3) - S$, so $\nabla(C(n, 3)) = \lceil \frac{n+1}{3} \rceil$ when $n = 3k$.

Case 2 $n = 3k + 1$ where k is a positive integer. Suppose that $S = \{a_t, n - 1\}$ where $a_1 = 1$, $a_t < n$ and

$$a_{t+1} = \begin{cases} a_t + 2, & \text{if } t \text{ is odd,} \\ a_t + 4, & \text{if } t \text{ is even.} \end{cases}$$

The graph $C(3k + 1, 3) - \{a_t\}$ when k is odd is acyclic, so is graph $C(3k + 1, 3) - S$ when k is even. Summarizing above, it is obtained that $\nabla(C(n, 3)) = \lceil \frac{n+1}{3} \rceil$ when $n = 3k + 1$.

Case 3 $n = 3k + 2$ where k is a positive even integer. Suppose that $S = \{a_t\}$ where $a_1 = 1, a_t < n$ and

$$a_{t+1} = \begin{cases} a_t + 2, & \text{if } t \text{ is odd,} \\ a_t + 4, & \text{if } t \text{ is even.} \end{cases}$$

Graph $C(3k+2) - S$ is acyclic. So $\nabla(C(n,3)) = \lceil \frac{n+1}{3} \rceil$ when $n = 3k + 2$.

Case 4 $n = 3k + 2$ where k is a positive odd integer. It can be seen that $n = 6 \times \frac{k-1}{2} + 5 = 6m + 5$. We prove that $\nabla(C(n,3)) \neq \lceil \frac{n+1}{3} \rceil = 2m + 2$. By contradiction. Suppose S is a decycling set of $C(n,3)$ and $|S| = 2m + 2$. Make the numbers in S with ascending order, that is, $S = \{b_1, b_2, \dots, b_{|S|}\}$ where $b_1 < b_2 < \dots < b_{|S|}$. Since S is a decycling set, then $|E(S)| = 8m + 8$. We know S is an independent set. For any two vertices b_i and b_{i+1} in S , $\chi(b_{i+1} - b_i) \neq 1$ or 3 and $\chi(b_{i+1} - b_i) \leq 4$. Suppose there are x pairs of numbers $(b_i, b_{i+1}) \in S$ such that $\chi(b_{i+1} - b_i) = 2$ and y pairs of numbers $(b_i, b_{i+1}) \in S$ such that $\chi(b_{i+1} - b_i) = 4$. Then two equation are derived: $x + y = 2m + 2$ (the number of the vertices whose removal is necessary) and $x + 3y = 4m + 3$. It follows that $2y = 2m + 1$. This is impossible. So $\nabla(C(n,3)) \neq \lceil \frac{n+1}{3} \rceil = 2m + 2$. That is $\nabla(C(6m+5,3)) \geq 2m + 3$. Let $S = \{a_t, n-1\}$ where $a_1 = 1, a_t < n$ and

$$a_{t+1} = \begin{cases} a_t + 2, & \text{if } t \text{ is odd,} \\ a_t + 4, & \text{if } t \text{ is even.} \end{cases}$$

The graph $C(3k,2) - S$ is acyclic when k is a positive odd integer, so $\nabla(C(3k+2,3)) = \lceil \frac{3k+3}{3} \rceil + 1 = k + 2$.

This theorem is found. \square

To obtain more results on decycling number, we introduce two operators. They are used in the following proof.

Given a labeled circular graph G with order n , suppose S is a set of vertices $S = \{i_1, i_2, \dots, i_l\}$. Define two operators δ and δ' such that $\delta S = \{i_j - i_{j-1} - 1, i_1 - i_l - 1 \mid j = 2, 3, \dots, l\}$ and $\delta' S = \{i_j - i_{j-1} - 1 \mid j = 2, 3, \dots, l\}$ where the elements read modulo n . From sets $\delta S = \{p_1, p_2, \dots, p_l\}$ or $\delta' S = \{p_1, p_2, \dots, p_{l-1}\}$ it is able to get the set S too. But the set S is different according to the choice of number i_1 .

For example, suppose G is the circular graph $C(17,8)$ and $S = \{3, 6, 7, 9, 10, 15, 17\}$. Then $\delta S = \{2, 0, 1, 0, 4, 1, 2\}$ and $\delta' S = \{2, 0, 1, 0, 4, 1\}$. If $\delta S = \{2, 1, 1, 0, 3, 1, 2\}$ and suppose $i_1 = 1$ the set S could be $\{1, 4, 6, 8, 9, 13, 15\}$.

Theorem 8

$$\nabla(C(n,4)) = \begin{cases} \lceil \frac{n+1}{3} \rceil + 1, & n = 3k + 2 \text{ and } k \text{ is a positive integer,} \\ \lceil \frac{n+1}{3} \rceil, & \text{otherwise.} \end{cases}$$

where $n \geq 9$.

Proof From Theorem 1, we know that $\nabla(C(n,4)) \geq \lceil \frac{n+1}{3} \rceil$. To get the lower bound, three cases are divided.

Case 1 $n = 3k$ where k is a positive integer. Let $S_0 = \{3i \mid i = 1, 2, \dots, k\}$. In graph $C(3k,4) - S_0$, for the vertex $v = 3i + 1$, two edges $(3i + 1, 3i)$ and $(3i + 1, 3i - 3)$ are removed. For the vertex $v = 3i + 2$, the edges $(3i + 2, 3i + 3)$ and $(3i + 2, 3i + 6)$ are removed. So the graph $C(3k,4) - S_0$ is a cycle of length $2k$. Removing any vertex of graph

$C(3k, 4) - S_0$, without loss of generality, suppose the vertex 1 is removed, the resultant graph is acyclic. So $S = S_0 \cup \{1\}$ is a ∇ -set. Then $\nabla(C(3k, 4)) = k + 1 = \lceil \frac{n+1}{3} \rceil$.

Case 2 $n = 3k + 1$ where k is a positive integer. Let $S = \{3i \mid i = 1, 2, \dots, k\} \cup \{1\}$. In graph $C(3k, 4) - S$, for the vertex $v = 3i + 1$ ($1 \leq i < k - 1$), it is incident to two edges $(3i + 1, 3i + 2)$ and $(3i + 1, 3i + 5)$. For the vertex $v = 3i + 2$ ($i \geq 2$), it is incident to two edges $(3i + 2, 3i + 1)$ and $(3i + 2, 3(i - 1) + 1)$. The vertices $2, n, 5, n - 3$ are articulated and the edges that they are incident with are $(2, n - 2)$, $(n, 4)$, $(5, 1)$ and $(n - 3, 1)$ respectively. So the graph $C(3k, 4) - S$ is acyclic. Then $\nabla(C(3k + 1, 4)) = k + 1 = \lceil \frac{n+1}{3} \rceil$.

Case 3 $n = 3k + 2$ where k is a positive integer. First we say that $\nabla(C(3k + 2, 4)) \neq k + 1$. Otherwise suppose that S is a decycling set of $k + 1$ vertices. From Corollary 2, $|E(S)| = 4k + 4$. Assume the numbers in S are sorted with ascending order. For any vertex a and its successor b in S , $\chi(a - b) = 2$, $\chi(a - b) = 3$ or $\chi(a - b) = 5$. Suppose there are x pairs of numbers (a, b) such that $\chi(a - b) = 2$, y pairs of numbers (c, d) such that $\chi(c - d) = 3$ and z pairs of numbers (e, f) such that $\chi(e - f) = 5$. Two equations are followed that $x + y + z = k + 1$ and $x + 2y + 4z = 2k + 1$.

If $z = 0$ then $y = k$ and $x = 1$. Suppose $i, i + 2$ are two numbers of S . Then the vertex $i + 1$ in graph $C(n, 4) - S$ is an isolated vertex which contradicts Corollary 2.

Then we have $z \neq 0$. Suppose the vertices $i, i + 5 \in S$ and the vertices $i + 1, i + 2, i + 3, i + 4 \notin S$. One can see that $i + 6$ and $i + 9 \notin S$ since set S is an independent set. The vertex $i + 7 \in S$ because set S is a decycling set. And it forces vertex $i + 8 \notin S$. The vertex $i + 10 \in S$ otherwise a circuit $\{i + 2, i + 3, i + 4, i + 8, i + 9, i + 10, i + 6\}$ appears which contradicts that set S is a decycling set. Then the vertices $i + 11$ and $i + 14 \notin S$. In order to get an acyclic graph, one of the vertices $i + 12, i + 13$ should be removed.

If the vertex $i + 12 \in S$ one can see that the vertices $i + 13, i + 14, i + 15, i + 16 \notin S$. Another pair of numbers $a, b \in S$ such that $\chi(a - b) = 5$ are obtained. That is to say a sequence T_1 with $\delta' T_1 = \{4, 1, 2, 1\}$ is available.

If the vertex $i + 13 \in S$ one can obtain a sequence T_2 with $\delta' T_2 = \{4, 1, 2, 2, 2, \dots, 1\}$.

The difference of T_1 and T_2 lies in the number of pairs (a, b) with $\chi(a - b) = 2$. Whatever the sequence in set S is, the order of the graph is a multiple of 3 which contradicts that $n = 3k + 2$.

In the proof of the case $z \neq 0$, $n \geq 12$. If $n = 11$, we know $z \neq 0$. Otherwise suppose $(1, 6) \in S$, a cycle $\{3, 4, 11, 10\}$ occurs.

Summarizing above one know that $\nabla(C(3k + 2, 4)) \geq k + 2$. Suppose $S = \{3i \mid i = 1, 2, \dots, k\} \cup \{1, n\}$. The graph $C(n, 4) - S$ is acyclic. Then $\nabla(C(3k + 2, 4)) = k + 2$. \square

Theorem 9 Suppose $n = 3k$, $l = 3m - 1$ and $(k, m) = 1$ where $k \geq 3m$. Then $\nabla(C(n, l)) = k + 1 = \lceil \frac{n+1}{3} \rceil$.

Proof. At first let $S_0 = \{3i \mid i = 1, 2, \dots, k\}$. For any vertex $3i + 1$, $3i + 1 + l = 3(i + m) \in S_0$, then $d_{G-S_0}(3i + 1) \leq 2$. And $3i + 1 - l = 3(i - m) + 2 \notin S_0$, then $d_{G-S_0}(3i + 1) = 2$. For the vertex $3i + 2$, $d_{G-S_0}(3i + 2) = 2$ since $3i + 2 + l = 3(i + m) + 1 \notin S_0$ and $3i + 2 - l = 3(i - m + 1) \in S_0$ then $d_{G-S_0}(3i + 2) = 2$.

Now we prove that $G - S_0$ is a circuit but not the disjoint union of some circuits. It suffices to prove that the set $A = B = \{1, 4, 7, 10 \dots, n - 2\}$ where $A = \{1, 3m + 1, 6m + 1, 9m + 1 \dots\}$ and the elements are modulo n .

In set A , suppose there are two elements $k_1 \times 3m + 1 \equiv k_2 \times 3m + 1$. Without loss of generality we can assume that $k_1 > k_2$. There exists a positive number x such that

$k_1 \times 3m + 1 = k_2 \times 3m + 1 + x \times 3k$ which follows that $(k_1 - k_2)m = xk$. The left can be divided by m , so xk is a multiple of m . On the other hand, $1 \leq |k_1 - k_2| \leq k$ which follows that $\frac{xk}{m} \leq k$. Then we can get $x \leq m$ which contradicts that xk is a multiple of m . Then every two elements in A are different.

Let $S = S_0 \cup \{1\}$. The graph $G - S$ is a path of length $2k - 1$. □

Theorem 10 Suppose $n = 3k$, $l = 3m - 1$ and $(k, m) = 2$ where $k \geq 3m$. Then $\nabla(C(n, l)) = k + 1 = \lceil \frac{n+1}{3} \rceil$.

Proof. At first let $S_0 = \{3i \mid i = 1, 2, \dots, k\}$. For any vertex $3i + 1$, $3i + 1 + l = 3(i + m) \in S_0$, then $d_{G-S_0}(3i + 1) \leq 2$. And $3i + 1 - l = 3(i - m) + 2 \notin S_0$, then $d_{G-S_0}(3i + 1) = 2$. For the vertex $3i + 2$, $d_{G-S_0}(3i + 2) = 2$ since $3i + 2 + l = 3(i + m) + 1 \notin S_0$ and $3i + 2 - l = 3(i - m + 1) \in S_0$ then $d_{G-S_0}(3i + 2) = 2$. Removing the vertices in set S_0 , two circuits are obtained $C_1 = \{1, 2, 2 + l, 3 + l, \dots, n - l, n - l + 1\}$ and $C_2 = \{n - 1, n - 2, \dots, l - 1\}$. Let $S = S_0 - \{n\} + \{1, n - 1\}$. It can be verified that graph $C(n, l) - S$ is acyclic. So $\nabla(C(n, l)) = k + 1$. □

From these theorems we know the lower bound in Theorem 1 is best possible since it is reached.

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